

On the ‘Ordnungszahlen’ in Gentzen’s First Consistency Proof

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Introduction

- [1] Gentzen, Die Widerspruchsfreiheit der reinen Zahlentheorie. Math. Ann. 112 (1936).
- [2] Gentzen, Neue Fassung des Widerspruchsfreiheitsbeweises für die reine Zahlentheorie. Forsch. zur Logik u. zur Grundl. d. ex. Wiss. Neue Folge 4 (1938)
- [3] Kogan-Bernstein, Simplification of Gentzen’s reductions in the classical arithmetic. Zap.Nauch.Sem. LOMI 105 (1981)
English translation: J.MATH.SCI. Vol.22(3) (1983)

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- §1 Derivations [1]
- §2 Gentzen’s ‘Ordnungszahlen’ $(\mathcal{O}, <_{\mathbb{R}})$ [1]
- §3 Assignment of ‘Ordnungszahlen’ to derivations [1]
- §4 Embedding of $(\mathcal{O}, <_{\mathbb{R}})$ into $(On, <)$
- §5 Transforming Gentzen’s assignment into the assignment in [3]
- §6 Reduction steps on sequents [1]
- §7 Reduction steps on derivations [1]
- §8 Transition to multisuccedent seqents (LK) yielding an “amalgamation” of [1] and [2]

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§1 Derivations. Inductive Definition of $d \vdash \Pi$.

1. Atomic derivations (axioms)
2. $d_0 \vdash \Gamma \rightarrow F(a)$ & variable condition
 $\implies I_{\forall x F}^a d_0 \vdash \Gamma \rightarrow \forall x F(x)$.
3. $d_0 \vdash \Gamma, A \rightarrow \perp \implies I_{\neg A} d_0 \vdash \Gamma \rightarrow \neg A$.
4. $d_0 \vdash \Gamma \rightarrow F(0)$ & $d_1 \vdash F(a), \Delta \rightarrow F(Sa)$ & var.cond.
 $\implies \text{Ind}_F^{a,t} d_0 d_1 \vdash \Gamma, \Delta \rightarrow F(t)$
5. If $d_i \vdash \Pi_i$ ($i = 0, \dots, l$), and if $\frac{\Pi_0 \dots \Pi_l}{\Pi}$ is a chain inference of rank r , then $\mathbf{K}_{\Pi}^r d_0 \dots d_l \vdash \Pi$.

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§2 Gentzen’s ‘Ordnungszahlen’ $(\mathcal{O}, <_{\mathbb{R}})$

$u, v \in \{0, 1\}^+$, $0^n := \overbrace{0 \dots 0}^n$, $n.u \in \mathbb{R}$

Definition of $M_n \subseteq \{0, 1\}^+$ (Mantissen)

1. $M_0 := \{1\}$;
2. $M_{n+1} := \{u_0 0^{n+1} u_1 0^{n+1} \dots 0^{n+1} u_l : u_0, \dots, u_l \in M_n \text{ \& } 0.u_l \leq_{\mathbb{R}} \dots \leq_{\mathbb{R}} 0.u_0\}$.

$M := \bigcup_{n \in \mathbb{N}} M_n$, $h(u) := \min\{n : u \in M_n\}$

Remarks. (a) $M_n \subseteq M_{n+1}$.

(b) $h(u)$ is the maximal number of consecutive zeros in u (Höchstanzahl aufeinanderfolgender Nullen).

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Definition.

$\mathcal{O} := \{n.u : n \in \mathbb{N} \text{ \& } u \in M_n\}$ (Ordnungszahlen)

Theorem 1. \mathcal{O} is wellordered by $<_{\mathbb{R}}$.

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§3 Assignment of ‘Ordnungszahlen’ to derivations

Definition of $\rho(d) \in \mathbb{N}$ (Numerus) and $\mu(d) \in M$ (Mantisse)

Case $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l$ (w.l.o.g. $l > 0$).

Let σ be a permut. s.t. $0.\mu(d_{\sigma(0)}) \geq_{\mathbb{R}} \dots \geq_{\mathbb{R}} 0.\mu(d_{\sigma(l)})$.

$$\mu(d) := u_0 0^{\nu+1} u_1 0^{\nu+1} \dots 0^{\nu+1} u_l,$$

where $u_i := \mu(d_{\sigma(i)})$ and $\nu := \max\{\mathbf{h}(u_0), \dots, \mathbf{h}(u_l)\}$

Abbreviation. $\mathbf{h}(d) := \mathbf{h}(\mu(d))$, $\mathbf{h}'(d) := \mathbf{h}(d) - 1$.

Proposition. $\mathbf{h}'(d) = \max\{\mathbf{h}(d_0), \dots, \mathbf{h}(d_l)\}$ (***)

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The numerus $\rho(d)$ is defined by the equation
 $\rho(d) - \mathbf{h}(d) = \max(\{(\rho(d_i) - \mathbf{h}(d_i)) - 1 : i \leq l\} \cup \{r\})$.

Remark.

With $\mathbf{dg}(d) := \rho(d) - \mathbf{h}(d)$, this equation becomes

$$\mathbf{dg}(d) = \max(\{\mathbf{dg}(d_i) - 1 : i \leq l\} \cup \{r\}).$$

Definition.

$\text{Ord}(d) := \rho(d) \cdot \mu(d)$ (*Ordnungszahl* of d)

Since (by definition) $\mathbf{h}(\mu(d)) = \mathbf{h}(d) \leq \rho(d)$,

we have $\mu(d) \in M_{\rho(d)}$, i.e. $\text{Ord}(d) \in \mathcal{O}$.

Theorem 2. If d^- results from d by a *reduction step on derivations* then $\text{Ord}(d^-) <_{\mathbb{R}} \text{Ord}(d)$.

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§4 Embedding of $(\mathcal{O}, <_{\mathbb{R}})$ into $(On, <)$

Definition of $|u|_n \in On$ for $u \in M_n$

1. $|1|_0 := 0$.
2. $|u_0 0^{n+1} \dots 0^{n+1} u_l|_{n+1} := \omega^{|u_0|_n} + \dots + \omega^{|u_l|_n}$.

Lemma 1. For $u \in M_n$ the following holds:

- (a) $|u|_{n+k} = \omega_k(|u|_n)$,
- (b) $\omega_n(0) \leq |u|_n < \omega_{n+1}(0)$.

Definition. For $n.u \in \mathcal{O}$ let $|n.u| := |u|_n$.

Lemma 2.

$n.u, m.v \in \mathcal{O}$ & $n.u <_{\mathbb{R}} m.v \Rightarrow |n.u| < |m.v|$.

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Definition. $\mathbf{o}(d) = |\text{Ord}(d)|$.

Theorem 2 together with Lemma 2 yields: $\mathbf{o}(d^-) < \mathbf{o}(d)$.

Our goal is, to find a direct recursive definition of $\mathbf{o}(d)$ which does not use the assignment $\text{Ord}(d)$.

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§5 Transforming Gentzen’s assignment

Abbreviation. $\tilde{\mathbf{o}}(d) := |\mu(d)|_{\mathbf{h}(d)}$

Lemma 3. $\mathbf{o}(d) = \omega_{\mathbf{dg}(d)}(\tilde{\mathbf{o}}(d))$.

Proof:

$$\mathbf{o}(d) = |\rho(d) \cdot \mu(d)| = |\mu(d)|_{\rho(d)} \stackrel{\text{L.1a}}{=} \omega_{\rho(d) - \mathbf{h}(d)}(|\mu(d)|_{\mathbf{h}(d)}).$$

Lemma 4.

For $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l$ we have $\tilde{\mathbf{o}}(d) = \omega^{\alpha_0} \# \dots \# \omega^{\alpha_l}$
 with $\alpha_i := \omega_{\mathbf{h}'(d) - \mathbf{h}(d_i)}(\tilde{\mathbf{o}}(d_i))$

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Lemma 4.

For $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l$ we have $\tilde{\mathbf{o}}(d) = \omega^{\alpha_0} \# \dots \# \omega^{\alpha_l}$
 with $\alpha_i := \omega_{\mathbf{h}'(d) - \mathbf{h}(d_i)}(\tilde{\mathbf{o}}(d_i))$

Proof:

By definition, $\mu(d) = u_0 0^{\nu+1} \dots 0^{\nu+1} u_l$ with
 $u_{\sigma'(i)} = \mu(d_i)$ and $\nu = \max\{\mathbf{h}(u_0), \dots, \mathbf{h}(u_l)\}$.

Hence $\nu = \max\{\mathbf{h}(d_0), \dots, \mathbf{h}(d_l)\} = \mathbf{h}'(d) = \mathbf{h}(d) - 1$,

$$\tilde{\mathbf{o}}(d) \stackrel{\text{Def}}{=} |\mu(d)|_{\mathbf{h}(d)} = |\mu(d)|_{\nu+1} = \omega^{|u_0|_{\nu}} + \dots + \omega^{|u_l|_{\nu}},$$

$$|u_{\sigma'(i)}|_{\nu} = |\mu(d_i)|_{\nu} \stackrel{\text{L.1a}}{=} \omega_{\nu - \mathbf{h}(d_i)}(|\mu(d_i)|_{\mathbf{h}(d_i)}).$$

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Summary. For $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l$ we have

- $\text{dg}(d) = \max(\{\text{dg}(d_i) - 1 : i \leq l\} \cup \{r\})$
- $\text{h}(d) = \max\{\text{h}(d_0), \dots, \text{h}(d_l)\} + 1$
- $\tilde{\text{o}}(d) = \omega^{\alpha_0} \# \dots \# \omega^{\alpha_l}$ with $\alpha_i := \omega_{\text{h}'(d) - \text{h}(d_i)}(\tilde{\text{o}}(d_i))$
- $\text{o}(d) = \omega_{\text{dg}(d)}(\tilde{\text{o}}(d))$

Idea. Simplify this ordinal assignment by setting

$$\tilde{\text{o}}(d) := \omega^{\tilde{\text{o}}(d_0)} \# \dots \# \omega^{\tilde{\text{o}}(d_l)}$$

Then $\text{h}(d)$ becomes obsolete, and the proof of $\text{o}(d^-) < \text{o}(d)$ works as well with this simpler $\tilde{\text{o}}$.

Comparison with [3]: $\text{h}(d) = \text{excess of } d$, $\tilde{\text{o}}(d) = \text{FO}(d)$

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Preliminaries to part II (Gentzen 1936)

- Formulas A, B, C, F are build up from arithmetic prime formulas by \forall, \wedge, \neg .
In the following, \wedge is not mentioned.
- Sequents: $\Pi = \Gamma \rightarrow C$,
 $\text{L}(\Gamma \rightarrow C) := \Gamma$, $\text{R}(\Gamma \rightarrow C) := \{C\}$,
 $A, \Pi := A, \Gamma \rightarrow C = (\{A\} \cup \Gamma) \rightarrow C$,
 $\Pi, A := \Gamma \rightarrow A$!!!
- Π has **endform** $:\Leftrightarrow \top \in \text{R}(\Pi) \vee \perp \in \text{L}(\Pi)$.

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§6 Reduction steps on sequents

$$(\text{R}_{\forall x F}) \frac{\dots \Gamma \rightarrow F(n) \dots (n \in \mathbb{N})}{\Gamma \rightarrow \forall x F(x)} \quad (\text{'Wahlfreiheit'})$$

$$(\text{L}_{\forall x F}^k) \frac{F(k), \Gamma \rightarrow C}{\Gamma \rightarrow C} \quad \text{if } \forall x F \in \Gamma$$

$$(\text{R}_{\neg A}) \frac{A, \Gamma \rightarrow \perp}{\Gamma \rightarrow \neg A} \quad (\text{L}_{\neg A}^0) \frac{\Gamma \rightarrow A}{\Gamma \rightarrow C} \quad \text{if } \neg A \in \Gamma$$

$$\mathcal{I}(\Pi, n) := \begin{cases} \Pi, F(n) & \text{if } \mathcal{I} = \text{R}_{\forall x F} \\ F(k), \Pi & \text{if } \mathcal{I} = \text{L}_{\forall x F}^k \\ A, \Pi, \perp & \text{if } \mathcal{I} = \text{R}_{\neg A} \\ \Pi, A & \text{if } \mathcal{I} = \text{L}_{\neg A}^0 \end{cases}$$

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Definition of $\mathcal{I} \triangleright \Pi$ (\mathcal{I} is permissible for Π)

$$\text{R}_B \triangleright \Pi \quad :\Leftrightarrow B \in \text{R}(\Pi)$$

$$\text{L}_B^k \triangleright \Pi \quad :\Leftrightarrow B \in \text{L}(\Pi)$$

Gentzen's Kettenschluß

$$\frac{\Gamma_0 \rightarrow A_0 \quad \dots \quad \Gamma_l \rightarrow A_l}{\Gamma \rightarrow C}$$
 is a chain inference of degree r if

there exists a $j \leq l$ such that $A_j = C$ and

$$\forall i \leq j (\Gamma_i \subseteq \Gamma, A_0, \dots, A_{i-1}) \quad \text{and} \quad \forall i < j (\text{rk}(A_i) \leq r).$$

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§7 Reduction steps on derivations

For each derivation d , whose endsequent is not in endform we shall now define *the reduction step on d* and at the same time **prove** the following:

by such a step the derivation d is transformed into another derivation $d[n]$ and its endsequent Π is thereby modified in the following way:

At most one reduction step $\text{tp}(d)$ is carried out on the sequent. It may thus happen that an endsequent remains entirely unchanged.

In this case we set $\text{tp}(d) := \text{Rep}$ with $\text{Rep}(\Pi, n) := \Pi$.

Theorem 3. $d \vdash \Pi \Rightarrow d[n] \vdash \text{tp}(d)(\Pi, n)$.

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Inference symbols

R_B, L_B^k (reduction steps on sequents),

Rep, Ax_0 .

syntactic variable for inference symbols: \mathcal{I}

Definition.

$\text{Ax}_0 \triangleright \Pi \quad :\Leftrightarrow \Pi$ has endform

$\text{Rep} \triangleright \Pi \quad :\Leftrightarrow 0 = 0$.

$\text{Rep}(\Pi, n) := \Pi$

Remark. $\mathcal{I} \triangleright \emptyset \rightarrow \perp \Rightarrow \mathcal{I} = \text{Rep}$.

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Theorem 3. If $d \vdash \Pi$ then

- (i) $\text{tp}(d) \triangleright \Pi$;
- (ii) $\text{Ax}_0 \not\vdash \Pi \Rightarrow d[n] \vdash \text{tp}(d)(\Pi, n) \ (\forall n)$.

Corollary.

$$d \vdash \emptyset \rightarrow \perp \Rightarrow d[0] \vdash \emptyset \rightarrow \perp.$$

Proof:

$$\begin{aligned} d \vdash \emptyset \rightarrow \perp &\stackrel{(i)+\text{Remark}}{\Rightarrow} \\ d \vdash \emptyset \rightarrow \perp \ \&\ \text{tp}(d) = \text{Rep} &\stackrel{(ii)}{\Rightarrow} \\ d[0] \vdash \text{Rep}(\emptyset \rightarrow \perp, 0) &= \emptyset \rightarrow \perp. \end{aligned}$$

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Definition of $\text{tp}(d)$ and $d[n]$

Simultaneously one proves: $d \vdash \Pi \Rightarrow \text{tp}(d) \triangleright \Pi$.

5. $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l \vdash \Pi$ with $d_i \vdash \Pi_i \ (i \leq l)$.

Let $j_0 := \min\{j \leq l : \mathbf{R}(\Pi_j) = \mathbf{R}(\Pi)\}$.

5.1. d **critical**, i.e. $\forall i \leq j_0 (\text{tp}(d_i) \not\vdash \Pi)$.

By IH $\forall i \leq j_0 (\text{tp}(d_i) \triangleright \Pi_i)$. It follows that there exist $i < j \leq j_0$ and B, k s.t. $\text{tp}(d_i) = \mathbf{R}_B$ & $\text{tp}(d_j) = \mathbf{L}_B^k$.

$\text{tp}(d) := \text{Rep}$ and $d[0] := \mathbf{K}_{\Pi}^{r-1} d\{0\}d\{1\}$ with

$$\begin{aligned} d\{0\} &:= \mathbf{K}_{\text{tp}(d_i)(\Pi, k)}^r d_0 \dots d_{i-1} d_i[k] d_{i+1} \dots d_l, \\ d\{1\} &:= \mathbf{K}_{\text{tp}(d_j)(\Pi, 0)}^r d_0 \dots d_{j-1} d_j[0] d_{j+1} \dots d_l. \end{aligned}$$

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5.2. d **not critical**:

Let i be minimal such that $\text{tp}(d_i) \triangleright \Pi$.

5.2.1. d_i **critical**: $\text{tp}(d) := \text{Rep}$ and

$$d[0] := \mathbf{K}_{\Pi}^{r'} d_0 \dots d_{i-1} d_i\{0\} d_i\{1\} d_{i+1} \dots d_l$$

5.2.2. d_i **not critical**:

$\text{tp}(d) := \text{tp}(d_i)$,

$$d[n] := \mathbf{K}_{\text{tp}(d_i)(\Pi, n)}^r d_0 \dots d_{i-1} d_i[n] d_{i+1} \dots d_l.$$

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§8 Transition to multisuccedent seqents (LK)

$$\Pi = \Gamma \rightarrow \Delta,$$

$$\mathbf{L}(\Pi) := \Gamma, \mathbf{R}(\Pi) := \Delta.$$

$$A, \Pi := A, \Gamma \rightarrow \Delta := (\{A\} \cup \Gamma) \rightarrow \Delta \text{ (as before)}$$

$$\Pi, A := \Gamma \rightarrow \Delta, A := \Gamma \rightarrow (\Delta \cup \{A\})$$

Kettenschluß: $\frac{\Pi_0 \ \dots \ \Pi_l}{\Pi}$ is a chain inference of degree r if Π can be derived from Π_0, \dots, Π_l by a finite number of cuts of degree $\leq r$.

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Definition of $\text{tp}(d)$ and $d[n]$

5. $d = \mathbf{K}_{\Pi}^r d_0 \dots d_l \vdash \Pi$ with $d_i \vdash \Pi_i \ (i \leq l)$.

Let $j_0 := l$. Then literally as before.

Lemma (*Existence of a suitable cut*)

If $\frac{\Pi_0 \ \dots \ \Pi_l}{\Pi}$ is a chain inference of degree r , and if $\forall i \leq l (\mathcal{I}_i \triangleright \Pi_i \ \& \ \mathcal{I}_i \not\vdash \Pi)$ then

$\exists i, j \leq l \exists k, B (\mathcal{I}_i = \mathbf{R}_B \ \& \ \mathcal{I}_j = \mathbf{L}_B^k \ \& \ \text{rk}(B) \leq r)$.

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